

Papers by Bunch [6] and de Hoog [7] will give entry to the literature.

CITED REFERENCES AND FURTHER READING:

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- Forsythe, G.E., and Moler, C.B. 1967, *Computer Solution of Linear Algebraic Systems* (Englewood Cliffs, NJ: Prentice-Hall), §19. [1]
- Westlake, J.R. 1968, *A Handbook of Numerical Matrix Inversion and Solution of Linear Equations* (New York: Wiley). [2]
- von Mises, R. 1964, *Mathematical Theory of Probability and Statistics* (New York: Academic Press), pp. 394ff. [3]
- Levinson, N., Appendix B of N. Wiener, 1949, *Extrapolation, Interpolation and Smoothing of Stationary Time Series* (New York: Wiley). [4]
- Robinson, E.A., and Treitel, S. 1980, *Geophysical Signal Analysis* (Englewood Cliffs, NJ: Prentice-Hall), pp. 163ff. [5]
- Bunch, J.R. 1985, *SIAM Journal on Scientific and Statistical Computing*, vol. 6, pp. 349–364. [6]
- de Hoog, F. 1987, *Linear Algebra and Its Applications*, vol. 88/89, pp. 123–138. [7]

2.9 Cholesky Decomposition

If a square matrix \mathbf{A} happens to be symmetric and positive definite, then it has a special, more efficient, triangular decomposition. *Symmetric* means that $a_{ij} = a_{ji}$ for $i, j = 1, \dots, N$, while *positive definite* means that

$$\mathbf{v} \cdot \mathbf{A} \cdot \mathbf{v} > 0 \quad \text{for all vectors } \mathbf{v} \quad (2.9.1)$$

(In Chapter 11 we will see that positive definite has the equivalent interpretation that \mathbf{A} has all positive eigenvalues.) While symmetric, positive definite matrices are rather special, they occur quite frequently in some applications, so their special factorization, called *Cholesky decomposition*, is good to know about. When you can use it, Cholesky decomposition is about a factor of two faster than alternative methods for solving linear equations.

Instead of seeking arbitrary lower and upper triangular factors \mathbf{L} and \mathbf{U} , Cholesky decomposition constructs a lower triangular matrix \mathbf{L} whose transpose \mathbf{L}^T can itself serve as the upper triangular part. In other words we replace equation (2.3.1) by

$$\mathbf{L} \cdot \mathbf{L}^T = \mathbf{A} \quad (2.9.2)$$

This factorization is sometimes referred to as “taking the square root” of the matrix \mathbf{A} . The components of \mathbf{L}^T are of course related to those of \mathbf{L} by

$$L_{ij}^T = L_{ji} \quad (2.9.3)$$

Writing out equation (2.9.2) in components, one readily obtains the analogs of equations (2.3.12)–(2.3.13),

$$L_{ii} = \left(a_{ii} - \sum_{k=1}^{i-1} L_{ik}^2 \right)^{1/2} \quad (2.9.4)$$

and

$$L_{ji} = \frac{1}{L_{ii}} \left(a_{ij} - \sum_{k=1}^{i-1} L_{ik} L_{jk} \right) \quad j = i + 1, i + 2, \dots, N \quad (2.9.5)$$

If you apply equations (2.9.4) and (2.9.5) in the order $i = 1, 2, \dots, N$, you will see that the L 's that occur on the right-hand side are already determined by the time they are needed. Also, only components a_{ij} with $j \geq i$ are referenced. (Since \mathbf{A} is symmetric, these have complete information.) It is convenient, then, to have the factor \mathbf{L} overwrite the subdiagonal (lower triangular but not including the diagonal) part of \mathbf{A} , preserving the input upper triangular values of \mathbf{A} . Only one extra vector of length N is needed to store the diagonal part of \mathbf{L} . The operations count is $N^3/6$ executions of the inner loop (consisting of one multiply and one subtract), with also N square roots. As already mentioned, this is about a factor 2 better than LU decomposition of \mathbf{A} (where its symmetry would be ignored).

A straightforward implementation is

```
#include <math.h>

void choldc(float **a, int n, float p[])
Given a positive-definite symmetric matrix a[1..n][1..n], this routine constructs its Cholesky
decomposition,  $\mathbf{A} = \mathbf{L} \cdot \mathbf{L}^T$ . On input, only the upper triangle of a need be given; it is not
modified. The Cholesky factor  $\mathbf{L}$  is returned in the lower triangle of a, except for its diagonal
elements which are returned in p[1..n].
{
    void nrerror(char error_text[]);
    int i,j,k;
    float sum;

    for (i=1;i<=n;i++) {
        for (j=i;j<=n;j++) {
            for (sum=a[i][j],k=i-1;k>=1;k--) sum -= a[i][k]*a[j][k];
            if (i == j) {
                if (sum <= 0.0)          a, with rounding errors, is not positive definite.
                    nrerror("choldc failed");
                p[i]=sqrt(sum);
            } else a[j][i]=sum/p[i];
        }
    }
}
```

You might at this point wonder about pivoting. The pleasant answer is that Cholesky decomposition is extremely stable numerically, without any pivoting at all. Failure of choldc simply indicates that the matrix \mathbf{A} (or, with roundoff error, another very nearby matrix) is not positive definite. In fact, choldc is an efficient way to test *whether* a symmetric matrix is positive definite. (In this application, you will want to replace the call to nrerror with some less drastic signaling method.)

Once your matrix is decomposed, the triangular factor can be used to solve a linear equation by backsubstitution. The straightforward implementation of this is

```
void cholsl(float **a, int n, float p[], float b[], float x[])
Solves the set of n linear equations  $\mathbf{A} \cdot \mathbf{x} = \mathbf{b}$ , where a is a positive-definite symmetric matrix.
a[1..n][1..n] and p[1..n] are input as the output of the routine choldc. Only the lower
subdiagonal portion of a is accessed. b[1..n] is input as the right-hand side vector. The
solution vector is returned in x[1..n]. a, n, and p are not modified and can be left in place
for successive calls with different right-hand sides b. b is not modified unless you identify b and
x in the calling sequence, which is allowed.
{
    int i,k;
    float sum;

    for (i=1;i<=n;i++) {          Solve  $\mathbf{L} \cdot \mathbf{y} = \mathbf{b}$ , storing  $\mathbf{y}$  in  $\mathbf{x}$ .
        for (sum=b[i],k=i-1;k>=1;k--) sum -= a[i][k]*x[k];
        x[i]=sum/p[i];
    }
    for (i=n;i>=1;i--) {          Solve  $\mathbf{L}^T \cdot \mathbf{x} = \mathbf{y}$ .
        for (sum=x[i],k=i+1;k<=n;k++) sum += a[k][i]*x[k];
    }
}
```

```

    x[i]=sum/p[i];
  }
}

```

A typical use of `cholc` and `cholsl` is in the inversion of covariance matrices describing the fit of data to a model; see, e.g., §15.6. In this, and many other applications, one often needs \mathbf{L}^{-1} . The lower triangle of this matrix can be efficiently found from the output of `cholc`:

```

for (i=1;i<=n;i++) {
  a[i][i]=1.0/p[i];
  for (j=i+1;j<=n;j++) {
    sum=0.0;
    for (k=i;k<j;k++) sum -= a[j][k]*a[k][i];
    a[j][i]=sum/p[j];
  }
}

```

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- Gill, P.E., Murray, W., and Wright, M.H. 1991, *Numerical Linear Algebra and Optimization*, vol. 1 (Redwood City, CA: Addison-Wesley), §4.9.2.
- Dahlquist, G., and Bjorck, A. 1974, *Numerical Methods* (Englewood Cliffs, NJ: Prentice-Hall), §5.3.5.
- Golub, G.H., and Van Loan, C.F. 1989, *Matrix Computations*, 2nd ed. (Baltimore: Johns Hopkins University Press), §4.2.

2.10 QR Decomposition

There is another matrix factorization that is sometimes very useful, the so-called *QR decomposition*,

$$\mathbf{A} = \mathbf{Q} \cdot \mathbf{R} \quad (2.10.1)$$

Here \mathbf{R} is upper triangular, while \mathbf{Q} is orthogonal, that is,

$$\mathbf{Q}^T \cdot \mathbf{Q} = \mathbf{1} \quad (2.10.2)$$

where \mathbf{Q}^T is the transpose matrix of \mathbf{Q} . Although the decomposition exists for a general rectangular matrix, we shall restrict our treatment to the case when all the matrices are square, with dimensions $N \times N$.

Like the other matrix factorizations we have met (*LU*, *SVD*, *Cholesky*), *QR* decomposition can be used to solve systems of linear equations. To solve

$$\mathbf{A} \cdot \mathbf{x} = \mathbf{b} \quad (2.10.3)$$

first form $\mathbf{Q}^T \cdot \mathbf{b}$ and then solve

$$\mathbf{R} \cdot \mathbf{x} = \mathbf{Q}^T \cdot \mathbf{b} \quad (2.10.4)$$

by backsubstitution. Since *QR* decomposition involves about twice as many operations as *LU* decomposition, it is not used for typical systems of linear equations. However, we will meet special cases where *QR* is the method of choice.